

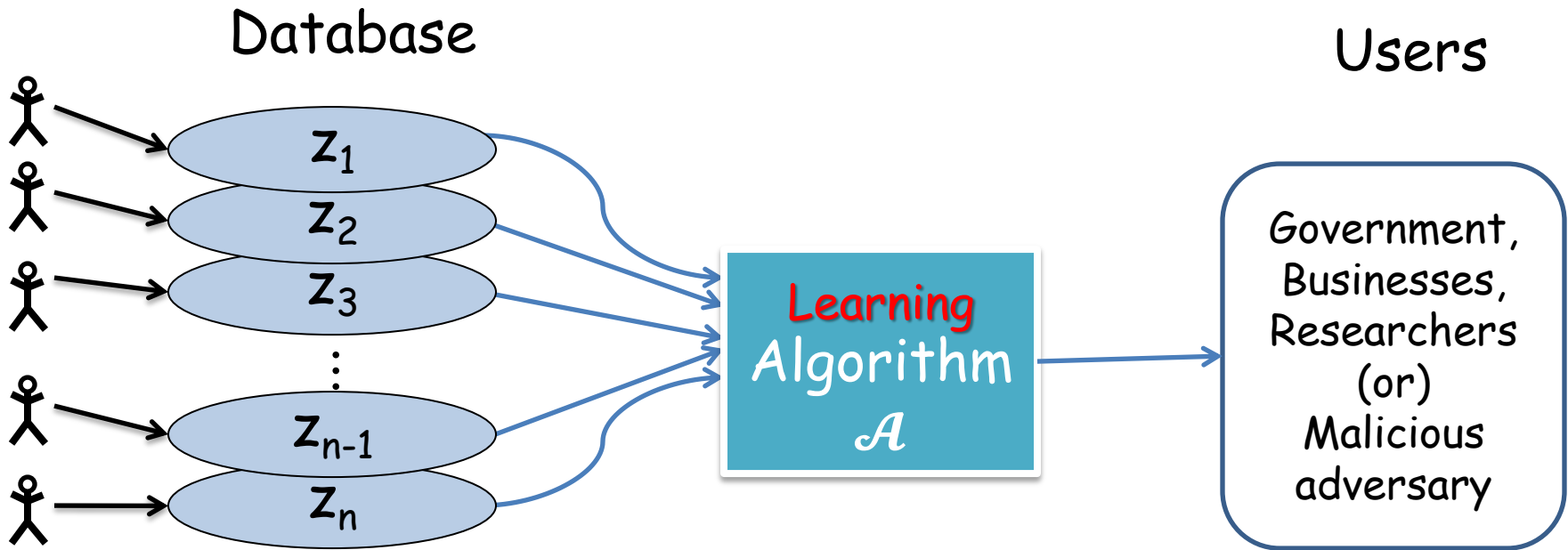
Characterizing the Sample Complexity of Private Learners

ITCS 2013

Uri Stemmer (BGU)

With Amos Beimel and Kobbi Nissim

Why Private Learners?



Often, this algorithmic task can be abstracted as a learning problem:

- Bank is interested in predicting (based on past customers) whether new customers are good/bad credit

Our Results:

- Introducing a new measure of concept classes:
Representation Dimension
- A combinatorial characterization for the **sample complexity** of private learners in terms of the representation dimension.
- Implications to **sanitization**, private optimization
 - Not in this talk

What is Private Learning?

Kasiviswanathan, Lee, Nissim, Raskhodnikova, Smith 08

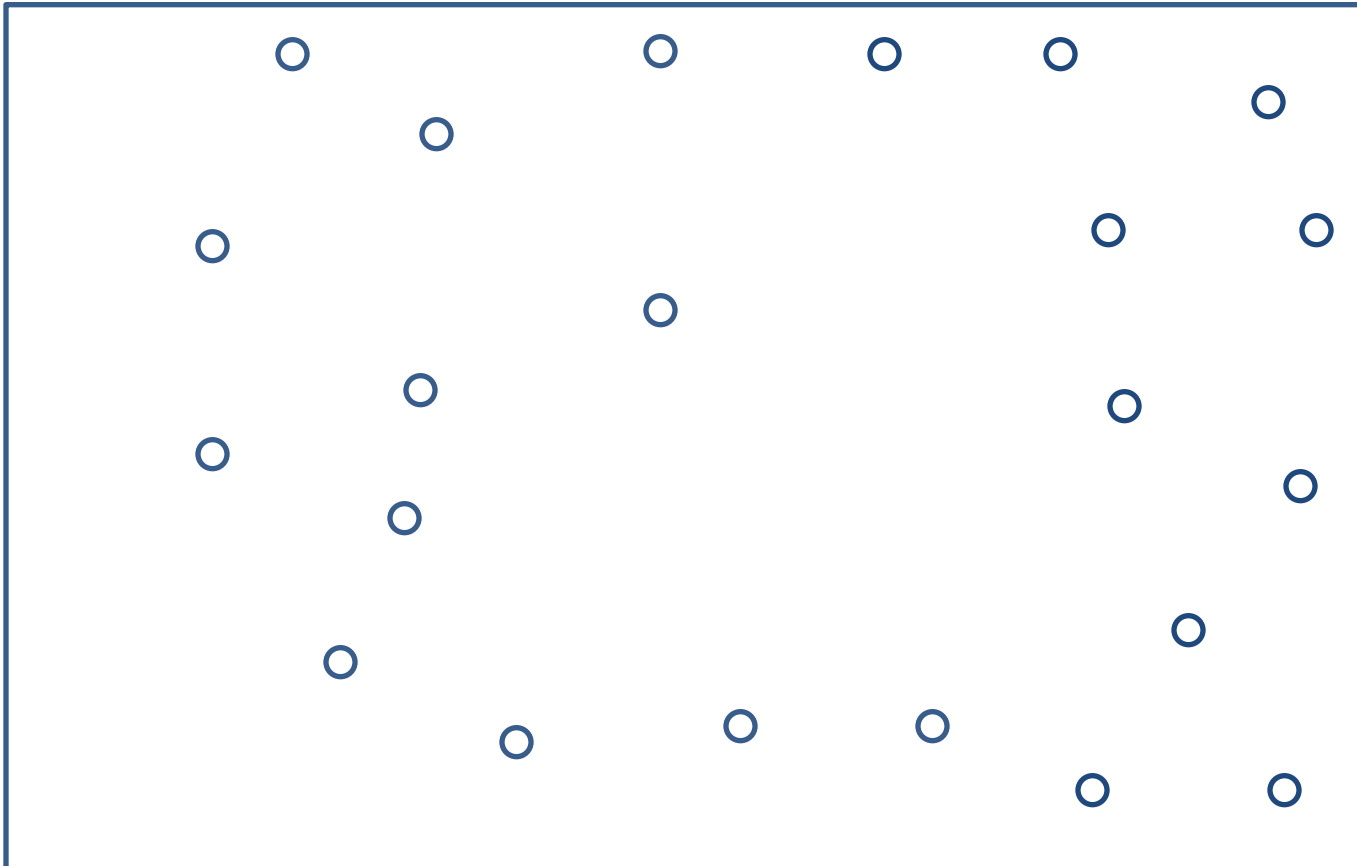
Definition:

+ PAC Learning
+ Differential Privacy

Private Learning

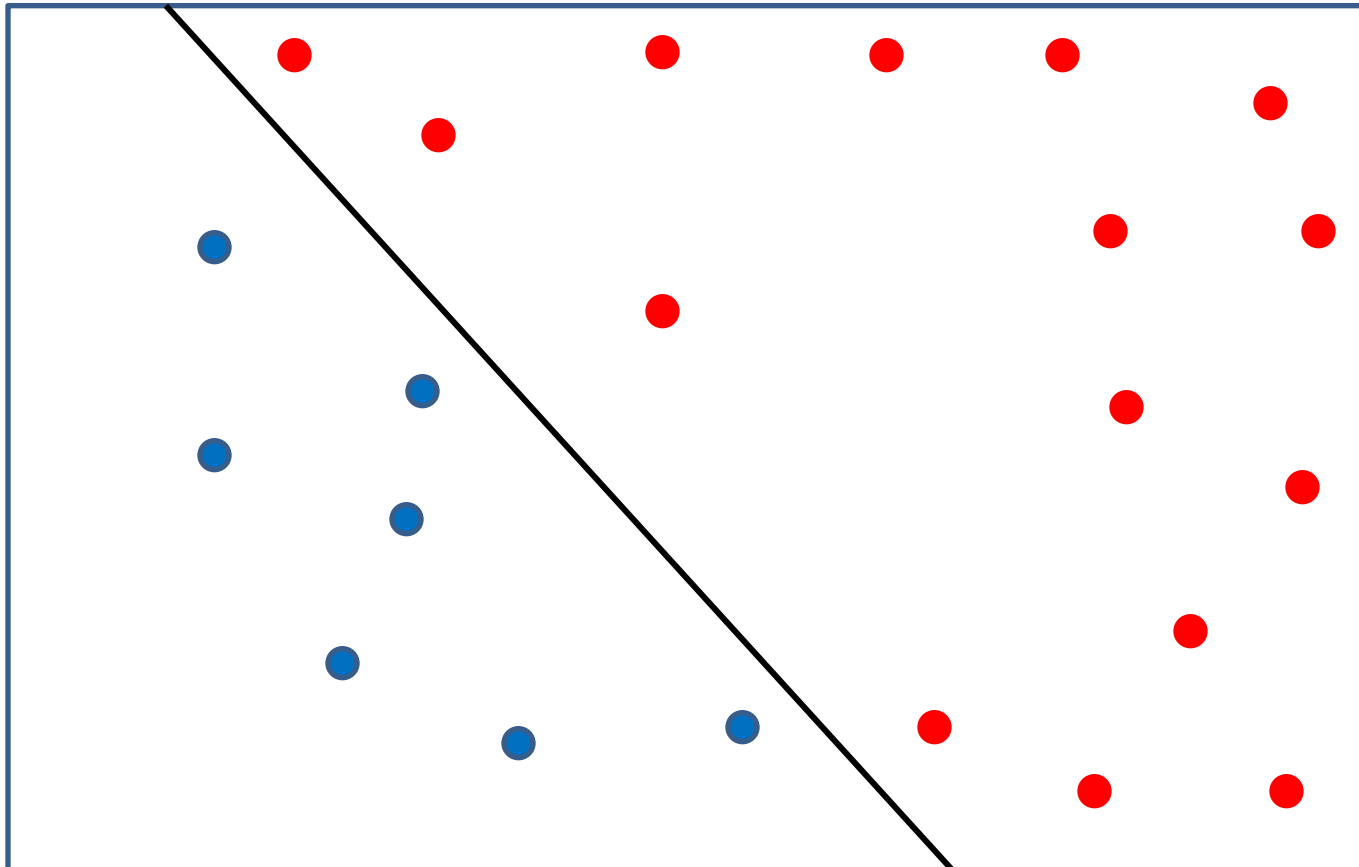
PAC Model [Valiant 84]

- A domain X
- \mathcal{C} = a set of boolean functions over X .



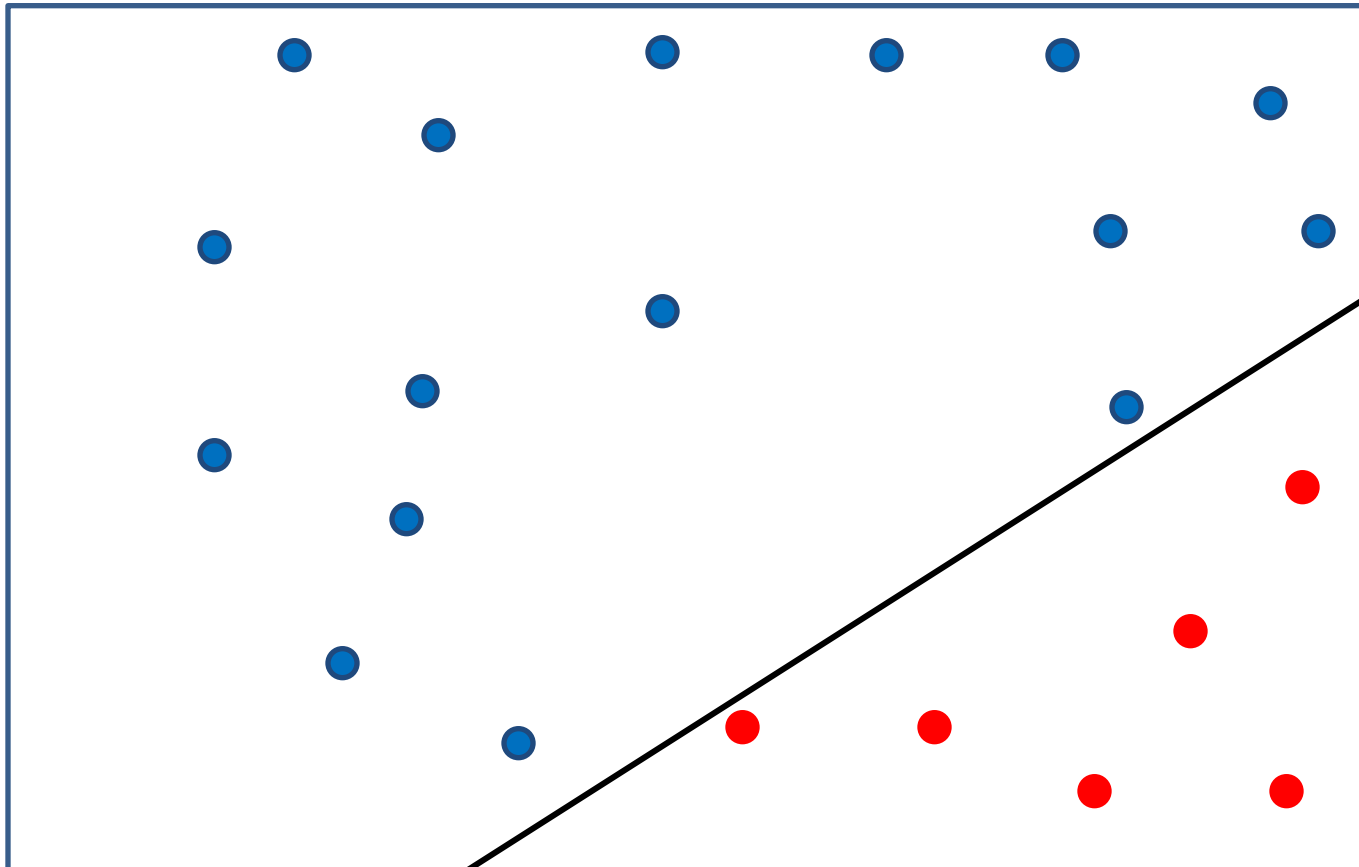
PAC Model [Valiant 84]

- A domain X
- \mathcal{C} = a set of boolean functions over X .



PAC Model [Valiant 84]

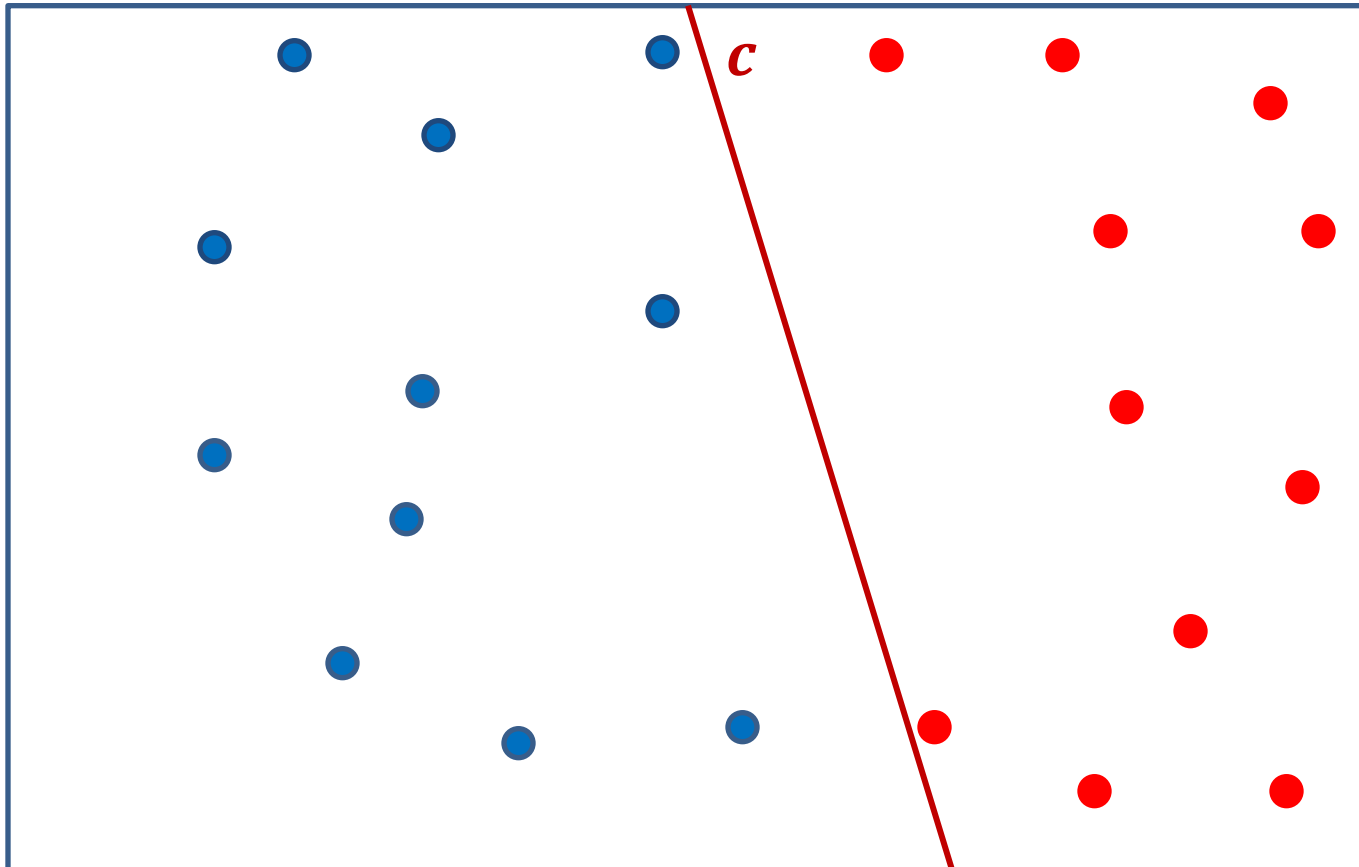
- A domain X
- \mathcal{C} = a set of boolean functions over X .



PAC Model [Valiant 84]

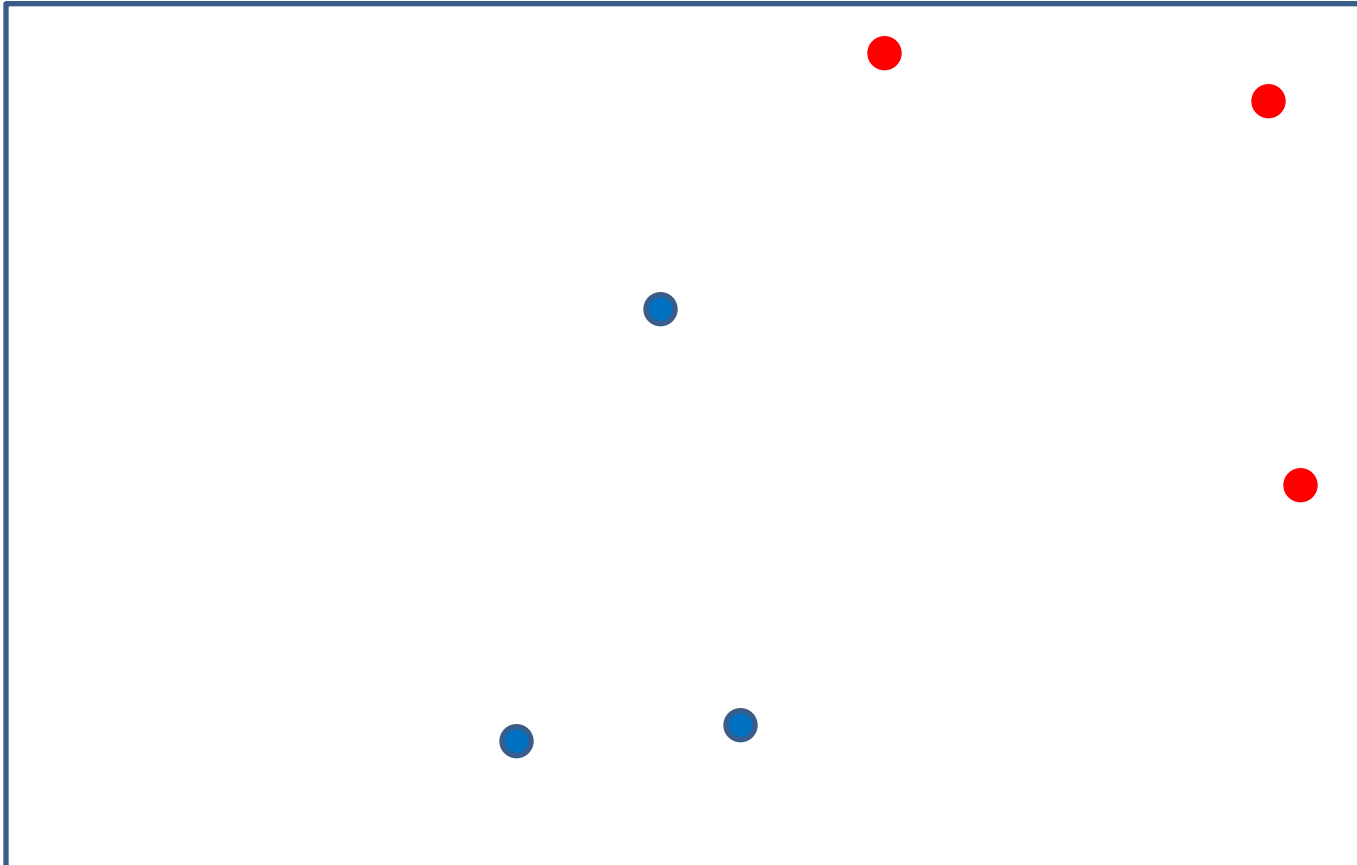
- A domain X
- \mathcal{C} = a set of boolean functions over X .

Unknown: - distribution \mathcal{D} over X .
- A specific target concept $c \in \mathcal{C}$.



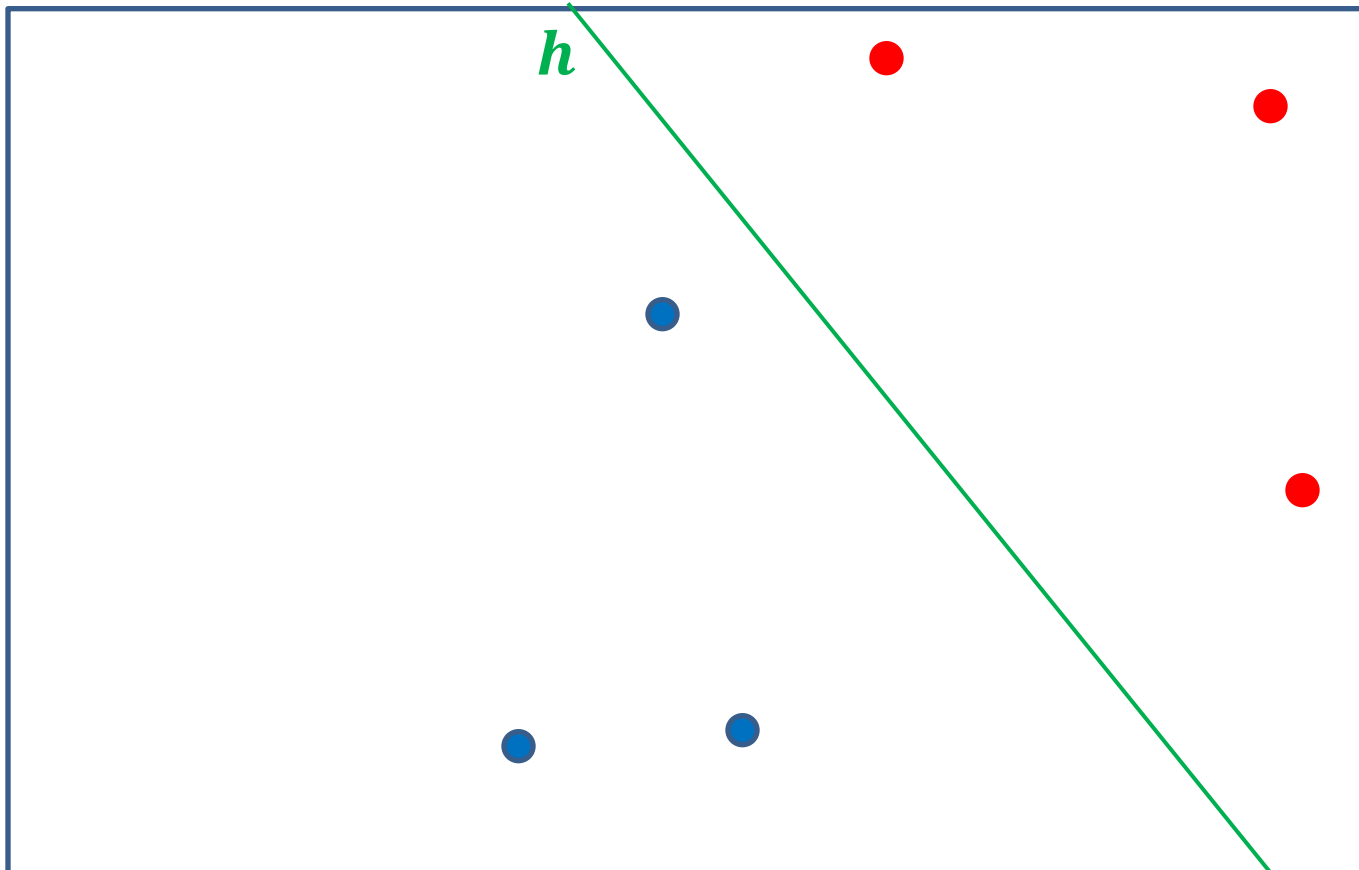
PAC Model [Valiant 84]

- Samples drawn according to \mathcal{D} and labeled by c



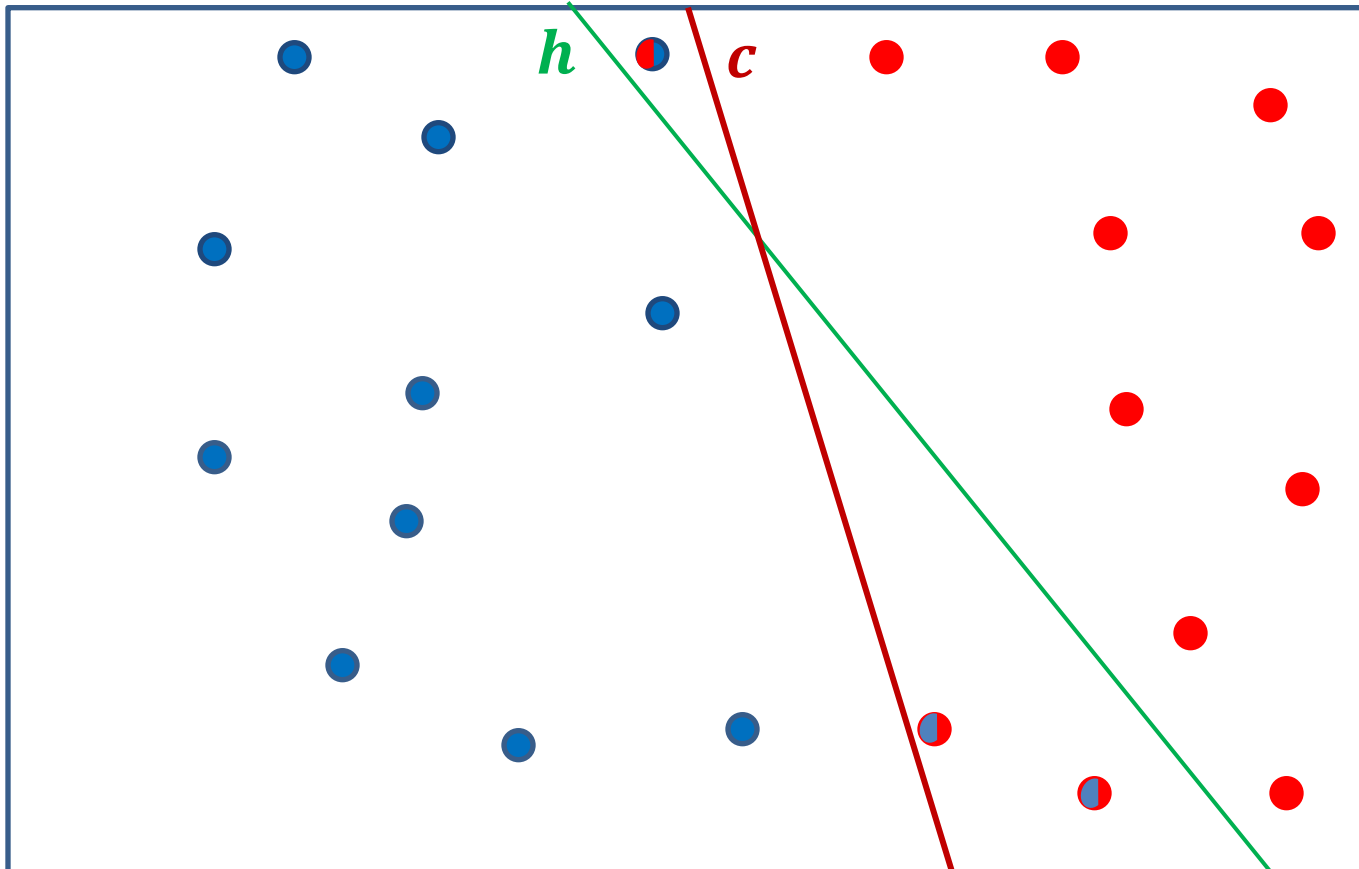
PAC Model [Valiant 84]

- Samples drawn according to \mathcal{D} and labeled by c
- A hypothesis h - a "guess" for c



PAC Model [Valiant 84]

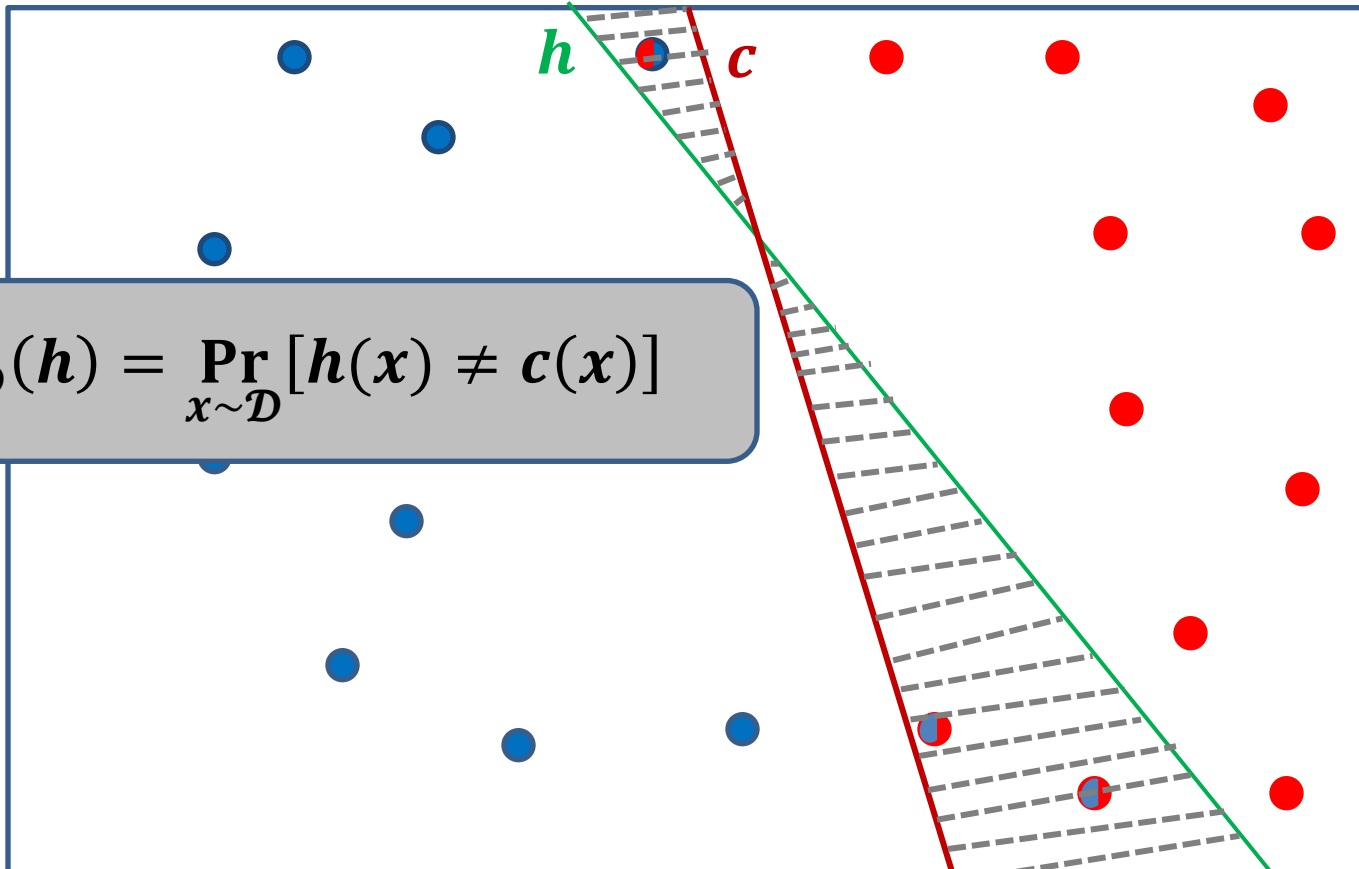
- Samples drawn according to \mathcal{D} and labeled by c
- A hypothesis h - a "guess" for c



PAC Model [Valiant 84]

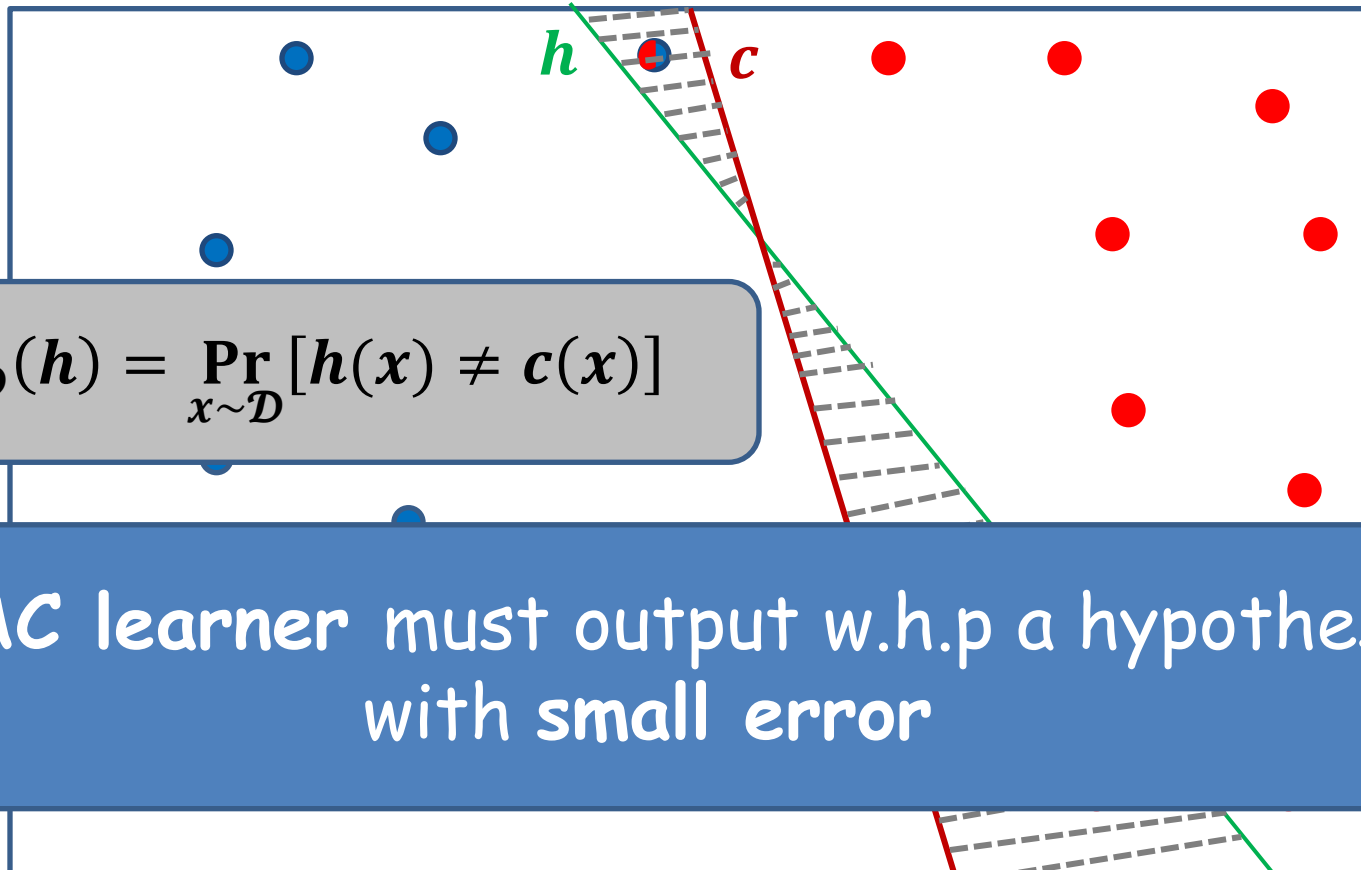
- Samples drawn according to \mathcal{D} and labeled by c
- A hypothesis h - a "guess" for c

$$\text{error}_{\mathcal{D}}(h) = \Pr_{x \sim \mathcal{D}} [h(x) \neq c(x)]$$



PAC Model [Valiant 84]

- Samples drawn according to \mathcal{D} and labeled by c
- A hypothesis h - a "guess" for c



$$\text{error}_{\mathcal{D}}(h) = \Pr_{x \sim \mathcal{D}} [h(x) \neq c(x)]$$

A PAC learner must output w.h.p a hypothesis with small error

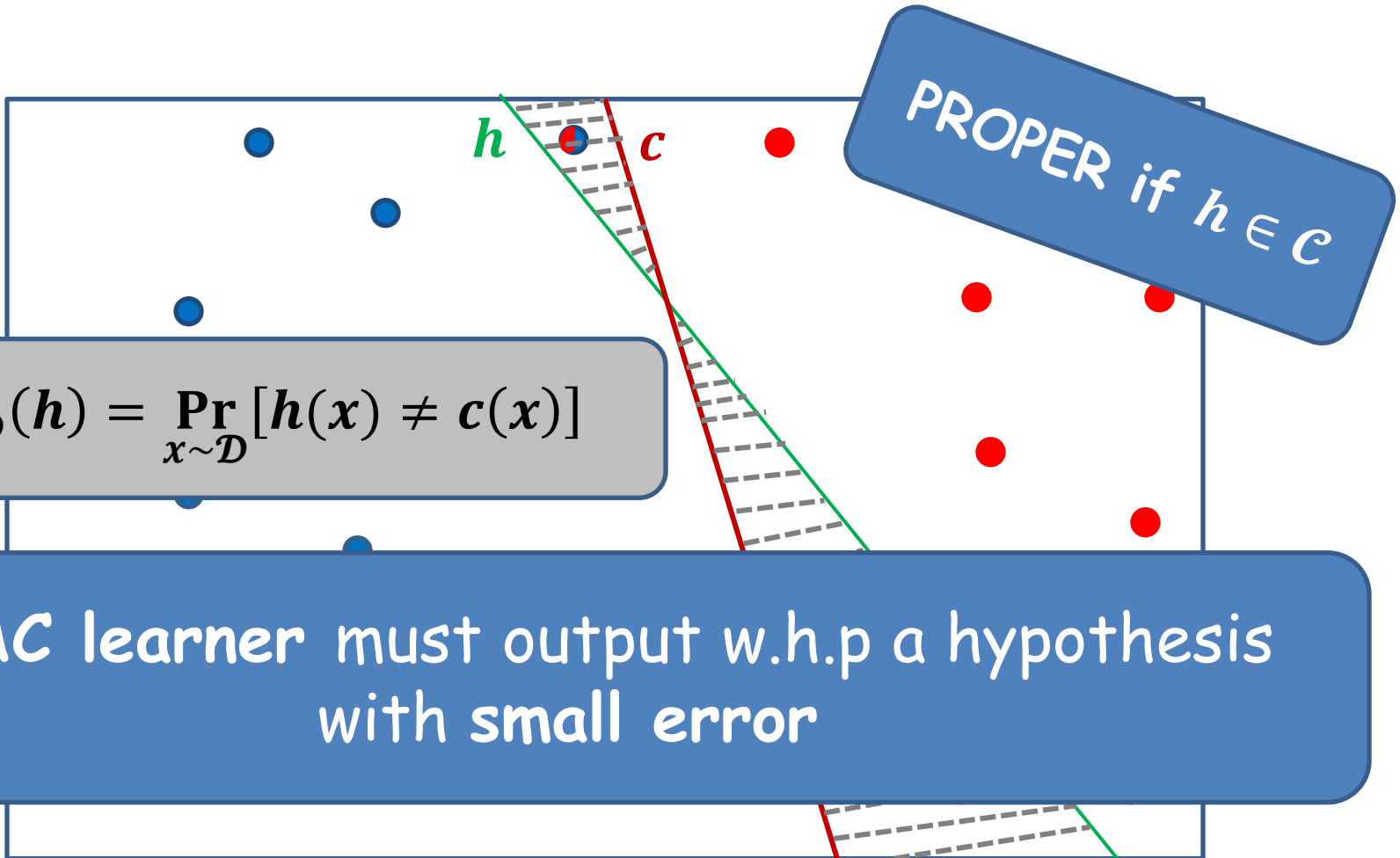
PAC Model [Valiant 84]

- Samples drawn according to \mathcal{D} and labeled by c
- A hypothesis h - a "guess" for c

$$\text{error}_{\mathcal{D}}(h) = \Pr_{x \sim \mathcal{D}} [h(x) \neq c(x)]$$

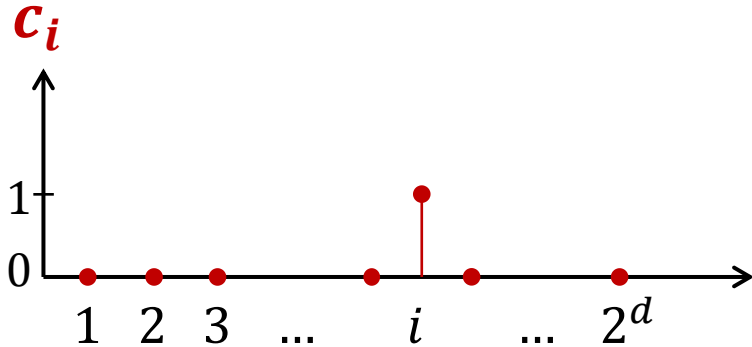
PROPER if $h \in \mathcal{C}$

A PAC learner must output w.h.p a hypothesis with small error



Example: POINT_d

Concept Class: $\mathcal{C} = \text{POINT}_d = \{c_1, \dots, c_{2^d}\}$



$$c_i(x) = 1 \iff x = i$$

A (non-private) learner for POINT_d with $O(1)$ samples:

- If there exists i s.t. $(i, 1)$ in the sample, return c_i .
- Otherwise (all labels are zero), return $h \equiv 0$.

What is Private Learning?

Kasiviswanathan, Lee, Nissim, Raskhodnikova, Smith 08

Definition:

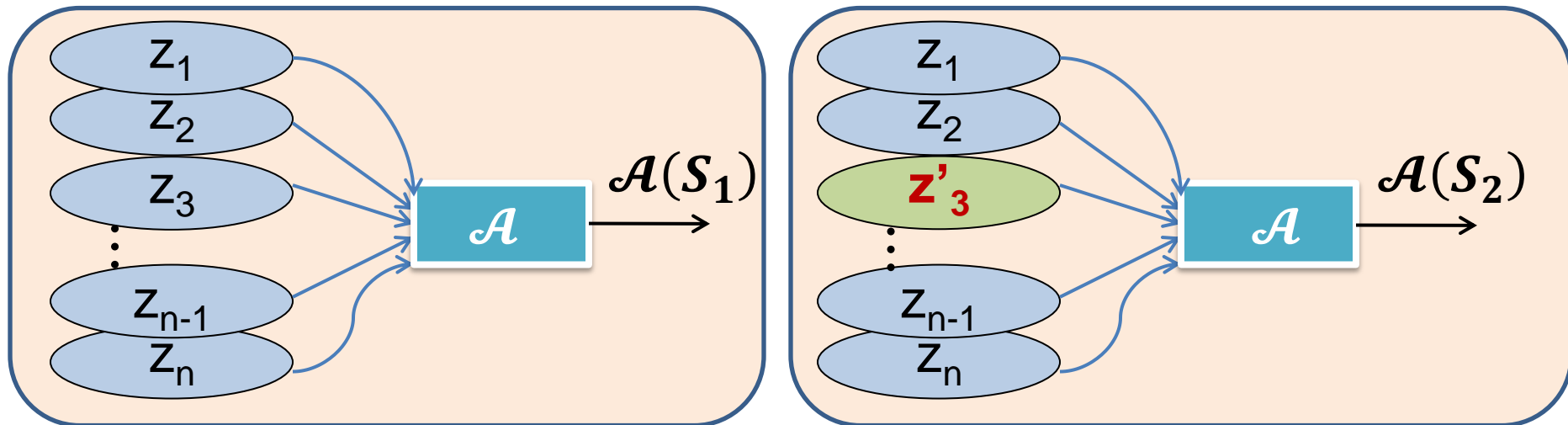
+ PAC Learning
+ Differential Privacy

Private Learning

Differential Privacy

Dwork, McSherry, Nissim, Smith 2006

Changing one record does not change the output distribution "too much"



Differential Privacy

Dwork, McSherry, Nissim, Smith 2006

Changing one record does not change the output distribution “too much”

A (rand) algorithm \mathcal{A} is differentially private if for all neighboring databases S_1, S_2 and for all sets of outputs F :

$$\Pr[\mathcal{A}(S_1) \in F] \approx \Pr[\mathcal{A}(S_2) \in F]$$

Differential Privacy

Dwork, McSherry, Nissim, Smith 2006

Changing one record does not change the output distribution “too much”

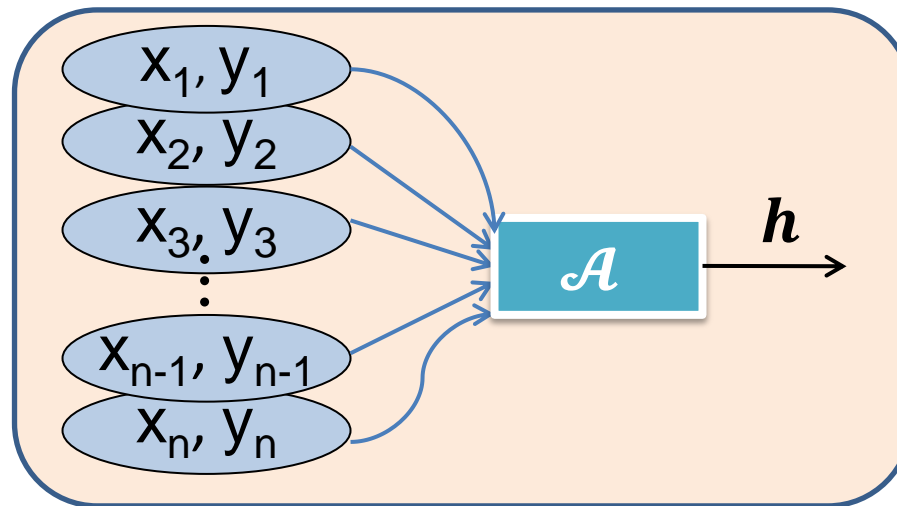
A (rand) algorithm \mathcal{A} is ϵ differentially private if for all neighboring databases S_1, S_2 and for all sets of outputs F :

$$\Pr[\mathcal{A}(S_1) \in F] \leq e^\epsilon \cdot \Pr[\mathcal{A}(S_2) \in F]$$

Private PAC (PPAC) [KLNRS 08]

Algorithm \mathcal{A} is a PPAC learner for a concept class \mathcal{C} if

- **Learning:** \mathcal{A} is a PAC learner for \mathcal{C} .
- **Privacy:** \mathcal{A} preserves differential privacy.



- Need to preserve privacy:
 - Label, sample, presence in database may all be sensitive!
- **Note:** privacy **not** preserved when we publish $h(x)$ for some x .

The **Sample** Complexity of Private Learners

Previous Results

Previous related work (partial list):

- [KLNRS08] Define private learning
- [BBKN10] Sample complexity of private learning
- [CH11] Learning in continuous domain, label privacy
- [CM08, CMS11, KST12] Machine learning
- [BLR08, DNRRV09, ...] Synthetic Data
- [Shapire90, ...] Boosting
- [DRV10] Private Boosting

Previous Work

Kasiviswanathan, Lee, Nissim, Raskhodnikova, Smith 08

Generic Construction (based on the Exp. Mechanism of [MT 07]):

Every finite concept class \mathcal{C} can be learned privately using $\log|\mathcal{C}|$ samples.

Previous Work

Kasiviswanathan, Lee, Nissim, Raskhodnikova, Smith 08

Generic Construction (based on the Exp. Mechanism of [MT 07]):

Every finite concept class \mathcal{C} can be learned privately using $\log|\mathcal{C}|$ samples.

- Recall: Learner for POINT_d with $O(1)$ samples (non-private).

- Generic construction of private learners results in $O(\log|\mathcal{C}|) = O(d)$ samples.

Is this gap essential?



Doom and Gloom?

[KLNRS 08]: A proper private PAC learner of POINT_d using $O(d)$ samples.

[BKN 10]: Proper private PAC learner of POINT_d must use $\Omega(d)$ samples.

[CH 11]: No proper PPAC learner of points in continuous domain.

Doom and Gloom?

[KLNRS 08]: A **proper** private PAC learner of POINT_d using $O(d)$ samples.

[BKN 10]: **Proper** private PAC learner of POINT_d must use $\Omega(d)$ samples.

[CH 11]: No proper PPAC learner of points in continuous domain.

Previous Work

Beimel, Brenner, Kasiviswanathan, Nissim 2010

Deterministic Representation: A hypothesis class \mathcal{H} s.t. for every $c \in \mathcal{C}$ and \mathcal{D} , there exists a hypothesis $h_0 \in \mathcal{H}$ s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.

Previous Work

Beimel, Brenner, Kasiviswanathan, Nissim 2010

Deterministic Representation: A hypothesis class \mathcal{H} s.t. for every $c \in \mathcal{C}$ and \mathcal{D} , there exists a hypothesis $h_0 \in \mathcal{H}$ s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.

Adversary

1. Choose $c \in \mathcal{C}$ and \mathcal{D}

c, \mathcal{D}

DRep \mathcal{H}

2. $\exists h_0 \in \mathcal{H}$ s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$

Previous Work

Beimel, Brenner, Kasiviswanathan, Nissim 2010

Deterministic Representation: A hypothesis class \mathcal{H} s.t. for every $c \in \mathcal{C}$ and \mathcal{D} , there exists a hypothesis $h_0 \in \mathcal{H}$ s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.

Adversary

1. Choose $c \in \mathcal{C}$ and \mathcal{D}

~~$c \mathcal{D}$~~

labeled sample

DRep \mathcal{H}

2. $\exists h_0 \in \mathcal{H}$ s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$

3. Choose $h \in \mathcal{H}$ using the Exp. Mechanism.

Reduces the sample complexity from $\log|\mathcal{C}|$ to $\log|\mathcal{H}|$

[BBKN 10]: $\log|\mathcal{H}|$ samples are NOT necessary.

Our Contribution:

Probabilistic Representation

A **list** of hypothesis classes $\{\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_r\}$ s.t. for every $c \in \mathcal{C}$ and \mathcal{D} :
w.p. $\frac{3}{4}$, **a randomly chosen \mathcal{H}_i** contains an h_0 s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.

Probabilistic Representation

A **list** of hypothesis classes $\{\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_r\}$ s.t. for every $c \in \mathcal{C}$ and \mathcal{D} :
w.p. $\frac{3}{4}$, **a randomly chosen \mathcal{H}_i** contains an h_0 s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.

Adversary

1. Choose $c \in \mathcal{C}$ and \mathcal{D}

c, \mathcal{D}



Prob. Rep.

1. Choose \mathcal{H}_i at random

2. **w.p. $\frac{3}{4}$** , $\exists h_0 \in \mathcal{H}_i$ s.t.
 $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$

Probabilistic Representation

A **list** of hypothesis classes $\{\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_r\}$ s.t. for every $c \in \mathcal{C}$ and \mathcal{D} :
w.p. $\frac{3}{4}$, **a randomly chosen \mathcal{H}_i** contains an h_0 s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.

Adversary

1. Choose $c \in \mathcal{C}$ and \mathcal{D}

~~c, \mathcal{D}~~ labeled sample \longrightarrow

Prob. Rep.

1. Choose \mathcal{H}_i at random
2. **w.p. $\frac{3}{4}$** , $\exists h_0 \in \mathcal{H}_i$ s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$
3. Choose $h \in \mathcal{H}_i$ using the Exp. Mechanism.

Reduces the sample complexity to **$\log|\mathcal{H}_i|$** .

RepDim

Probabilistic Representation:

A list of hypothesis classes $\{\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_r\}$ s.t. for every $c \in \mathcal{C}$ and \mathcal{D} :
w.p. $\frac{3}{4}$, a randomly chosen \mathcal{H}_i contains an h_0 s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.

RepDim

Probabilistic Representation:

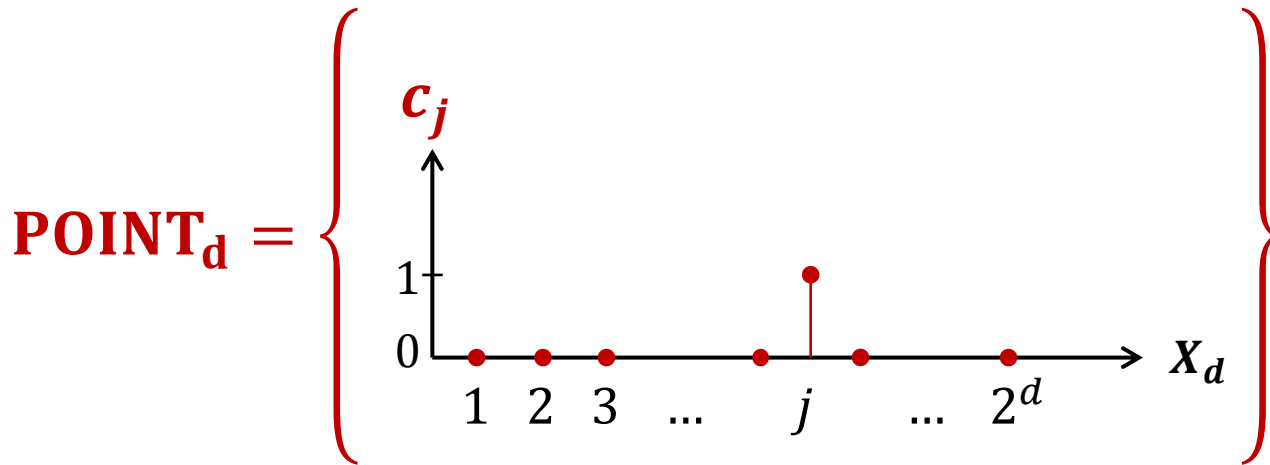
A list of hypothesis classes $\{\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_r\}$ s.t. for every $c \in \mathcal{C}$ and \mathcal{D} :
w.p. $\frac{3}{4}$, a randomly chosen \mathcal{H}_i contains an h_0 s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.

- The **size** of a prob. rep. is defined as $\max \log |\mathcal{H}_i|$.
- Define **RepDim**(\mathcal{C}) as the size of \mathcal{C} 's minimal probabilistic representation.

Example: Prob. Rep. for POINT_d

Probabilistic Representation:

A list of hypothesis classes $\{\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_r\}$ s.t. for every $c \in \mathcal{C}$ and \mathcal{D} :
w.p. $\frac{3}{4}$, a randomly chosen \mathcal{H}_i contains an h_0 s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.



$$\text{Rep} = \left\{ \mathcal{H} \subseteq 2^{X_d} : |\mathcal{H}| \leq 22 \right\}.$$

$$\text{Size}(\text{Rep}) = O(1).$$

Example: Prob. Rep. for POINT_d

$$\text{Rep} = \left\{ \mathcal{H} \subseteq 2^{X_d} : |\mathcal{H}| \leq 22 \right\}.$$

- **Choose \mathcal{H} as follows:**
 - For $i = 1 \dots 22$ construct h_i by setting $h_i(x)=1$ with probability $1/8$ and $h_i(x)=0$ otherwise.
 - $\mathcal{H} = \{h_i\}$.

Example: Prob. Rep. for POINT_d

$$\text{Rep} = \left\{ \mathcal{H} \subseteq 2^{X_d} : |\mathcal{H}| \leq 22 \right\}.$$

- **Choose \mathcal{H} as follows:**
 - For $i = 1 \dots 22$ construct h_i by setting $h_i(x)=1$ with probability $1/8$ and $h_i(x)=0$ otherwise.
 - $\mathcal{H} = \{h_i\}$.
- **Claim:** Fix a target concept $c_j \in \text{POINT}_d$ and a distribution \mathcal{D} . The probability that \mathcal{H} contains h_i s.t. $\text{error}_{\mathcal{D}}(h_i) \leq \frac{1}{4}$ is at least $\frac{3}{4}$.
- **Proof:**

Example: Prob. Rep. for POINT_d

$$\text{Rep} = \left\{ \mathcal{H} \subseteq 2^{X_d} : |\mathcal{H}| \leq 22 \right\}.$$

- **Choose \mathcal{H} as follows:**
 - For $i = 1 \dots 22$ construct h_i by setting $h_i(x)=1$ with probability $1/8$ and $h_i(x)=0$ otherwise.
 - $\mathcal{H} = \{h_i\}$.
- **Claim:** Fix a target concept $c_j \in \text{POINT}_d$ and a distribution \mathcal{D} . The probability that \mathcal{H} contains h_i s.t. $\text{error}_{\mathcal{D}}(h_i) \leq \frac{1}{4}$ is at least $\frac{3}{4}$.
- **Proof:**
 - $\mathbb{E}[\text{error}_{\mathcal{D}}(h_i) | h_i(j) = 1] \leq \frac{1}{8}$.

Example: Prob. Rep. for POINT_d

$$\text{Rep} = \left\{ \mathcal{H} \subseteq 2^{X_d} : |\mathcal{H}| \leq 22 \right\}.$$

- **Choose \mathcal{H} as follows:**
 - For $i = 1 \dots 22$ construct h_i by setting $h_i(x)=1$ with probability $1/8$ and $h_i(x)=0$ otherwise.
 - $\mathcal{H} = \{h_i\}$.
- **Claim:** Fix a target concept $c_j \in \text{POINT}_d$ and a distribution \mathcal{D} . The probability that \mathcal{H} contains h_i s.t. $\text{error}_{\mathcal{D}}(h_i) \leq \frac{1}{4}$ is at least $\frac{3}{4}$.
- **Proof:**
 - $\mathbb{E}[\text{error}_{\mathcal{D}}(h_i) | h_i(j) = \mathbf{1}] \leq \frac{1}{8}$.
 - By Markov's inequality $\Pr \left[\text{error}_{\mathcal{D}}(h_i) > \frac{1}{4} \mid h_i(j) = \mathbf{1} \right] \leq \frac{1}{2}$.

Example: Prob. Rep. for POINT_d

$$\text{Rep} = \left\{ \mathcal{H} \subseteq 2^{X_d} : |\mathcal{H}| \leq 22 \right\}.$$

- **Choose \mathcal{H} as follows:**

- For $i = 1 \dots 22$ construct h_i by setting $h_i(x)=1$ with probability $1/8$ and $h_i(x)=0$ otherwise.
- $\mathcal{H} = \{h_i\}$.

- **Claim:** Fix a target concept $c_j \in \text{POINT}_d$ and a distribution \mathcal{D} . The probability that \mathcal{H} contains h_i s.t. $\text{error}_{\mathcal{D}}(h_i) \leq \frac{1}{4}$ is at least $\frac{3}{4}$.

- **Proof:**

- $\mathbb{E}[\text{error}_{\mathcal{D}}(h_i) | h_i(j) = \mathbf{1}] \leq \frac{1}{8}$.
- By Markov's inequality $\Pr \left[\text{error}_{\mathcal{D}}(h_i) > \frac{1}{4} \mid h_i(j) = \mathbf{1} \right] \leq \frac{1}{2}$.
- $\Pr \left[\text{error}_{\mathcal{D}}(h_i) \leq \frac{1}{4} \right] \geq \Pr[h_i(j) = \mathbf{1}] \Pr \left[\text{error}_{\mathcal{D}}(h_i) \leq \frac{1}{4} \mid h_i(j) = \mathbf{1} \right] \geq \frac{1}{16}$.

Example: Prob. Rep. for POINT_d

$$\text{Rep} = \left\{ \mathcal{H} \subseteq 2^{X_d} : |\mathcal{H}| \leq 22 \right\}.$$

- **Choose \mathcal{H} as follows:**

- For $i = 1 \dots 22$ construct h_i by setting $h_i(x)=1$ with probability $1/8$ and $h_i(x)=0$ otherwise.
- $\mathcal{H} = \{h_i\}$.

- **Claim:** Fix a target concept $c_j \in \text{POINT}_d$ and a distribution \mathcal{D} . The probability that \mathcal{H} contains h_i s.t. $\text{error}_{\mathcal{D}}(h_i) \leq \frac{1}{4}$ is at least $\frac{3}{4}$.

- **Proof:**

- $\mathbb{E}[\text{error}_{\mathcal{D}}(h_i) | h_i(j) = \mathbf{1}] \leq \frac{1}{8}$.
- By Markov's inequality $\Pr \left[\text{error}_{\mathcal{D}}(h_i) > \frac{1}{4} \mid h_i(j) = \mathbf{1} \right] \leq \frac{1}{2}$.
- $\Pr \left[\text{error}_{\mathcal{D}}(h_i) \leq \frac{1}{4} \right] \geq \Pr[h_i(j) = \mathbf{1}] \Pr \left[\text{error}_{\mathcal{D}}(h_i) \leq \frac{1}{4} \mid h_i(j) = \mathbf{1} \right] \geq \frac{1}{16}$.
- \mathcal{H} fails to contains h_i s.t. $\text{error}_{\mathcal{D}}(h_i) \leq \frac{1}{4}$ w.p. at most $\left(1 - \frac{1}{16}\right)^{22} < \frac{1}{4}$.

Example: Prob. Rep. for POINT_d

$$\text{Rep} = \left\{ \mathcal{H} \subseteq 2^{X_d} : |\mathcal{H}| \leq 22 \right\}.$$

- **Choose \mathcal{H} as follows:**

- For $i = 1 \dots 22$ construct h_i by setting $h_i(x)=1$ with probability $1/8$ and $h_i(x)=0$ otherwise.
- $\mathcal{H} = \{h_i\}$.

- **Claim:** Fix a target concept $c_j \in \text{POINT}_d$ and a distribution \mathcal{D} . The probability that \mathcal{H} contains h_i s.t. $\text{error}_{\mathcal{D}}(h_i) \leq \frac{1}{4}$ is at least $\frac{3}{4}$.

- **Proof:**

- $\mathbb{E}[\text{error}_{\mathcal{D}}(h_i) | h_i(j) = \mathbf{1}] \leq \frac{1}{8}$.
- By Markov's inequality $\Pr \left[\text{error}_{\mathcal{D}}(h_i) > \frac{1}{4} \mid h_i(j) = \mathbf{1} \right] \leq \frac{1}{2}$.
- $\Pr \left[\text{error}_{\mathcal{D}}(h_i) \leq \frac{1}{4} \right] \geq \Pr[h_i(j) = \mathbf{1}] \Pr \left[\text{error}_{\mathcal{D}}(h_i) \leq \frac{1}{4} \mid h_i(j) = \mathbf{1} \right] \geq \frac{1}{16}$.
- \mathcal{H} fails to contain h_i s.t. $\text{error}_{\mathcal{D}}(h_i) \leq \frac{1}{4}$ w.p. at most $\left(1 - \frac{1}{16}\right)^{22} < \frac{1}{4}$.

- **Efficiency:** enough if entries of h_i are pairwise ind.

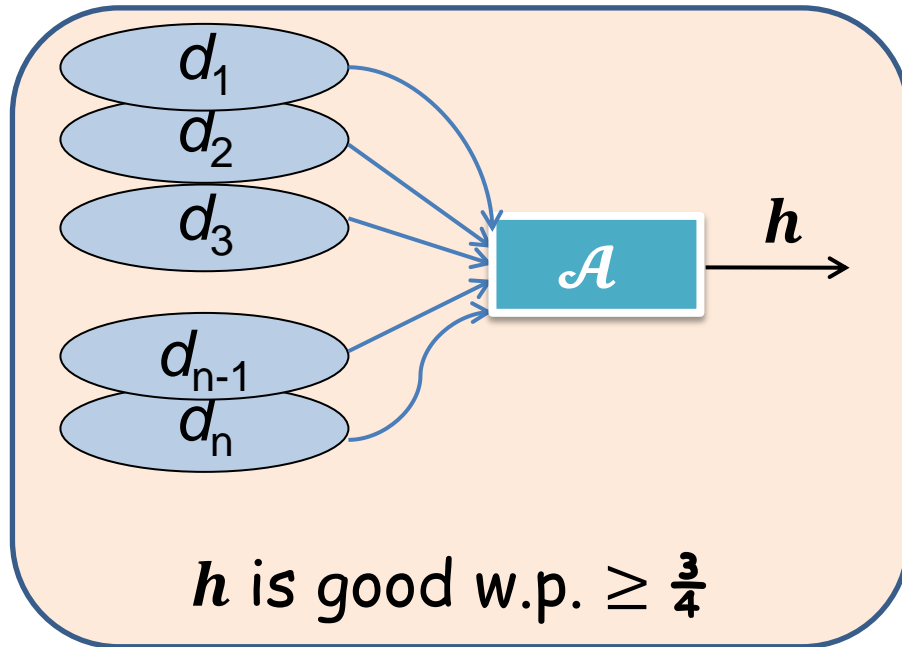
Until Now

- Defined **Probabilistic Representation**.
- We showed:
Probabilistic Representation \Rightarrow Private Learner.

Private Learner \implies Prob. Rep.

\mathcal{A} : PPAC learner for a concept class C using n examples.

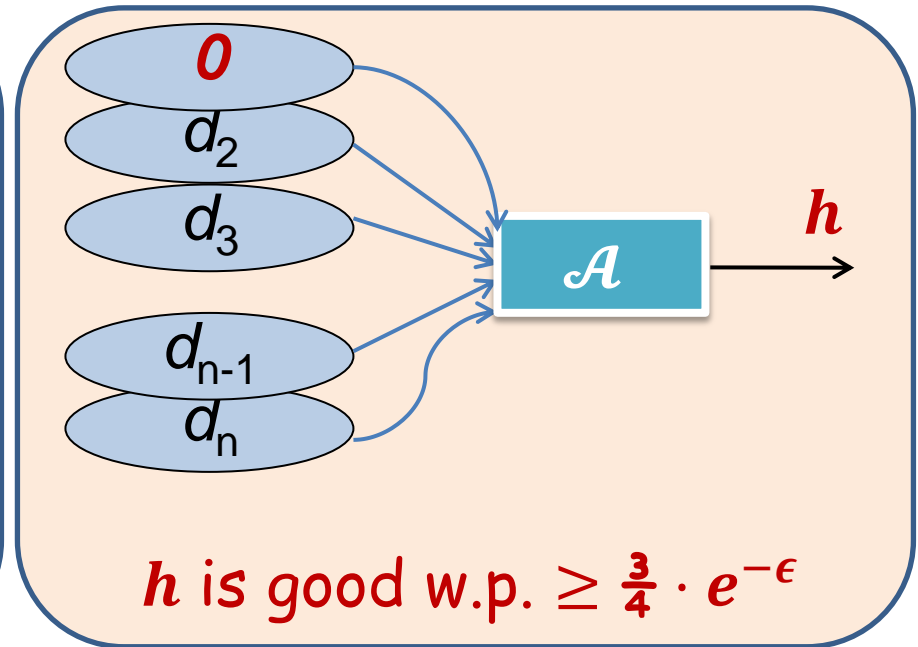
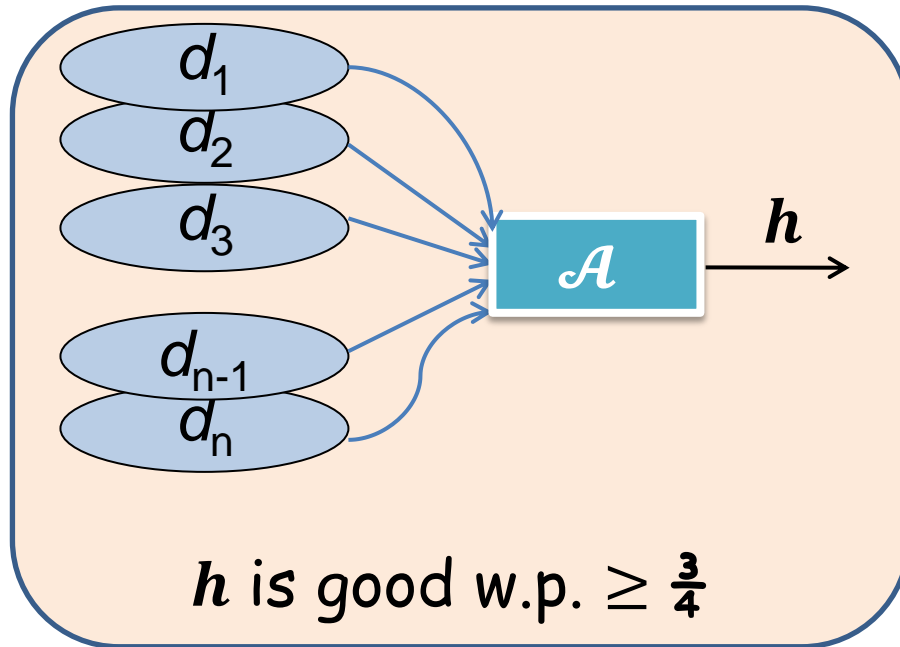
Correctly sampled data:



Private Learner \implies Prob. Rep.

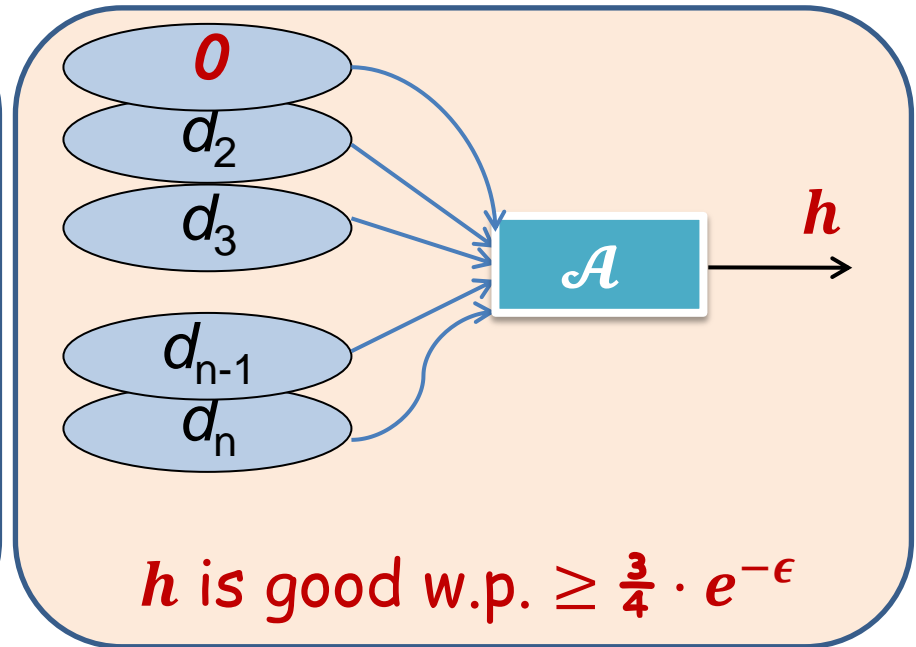
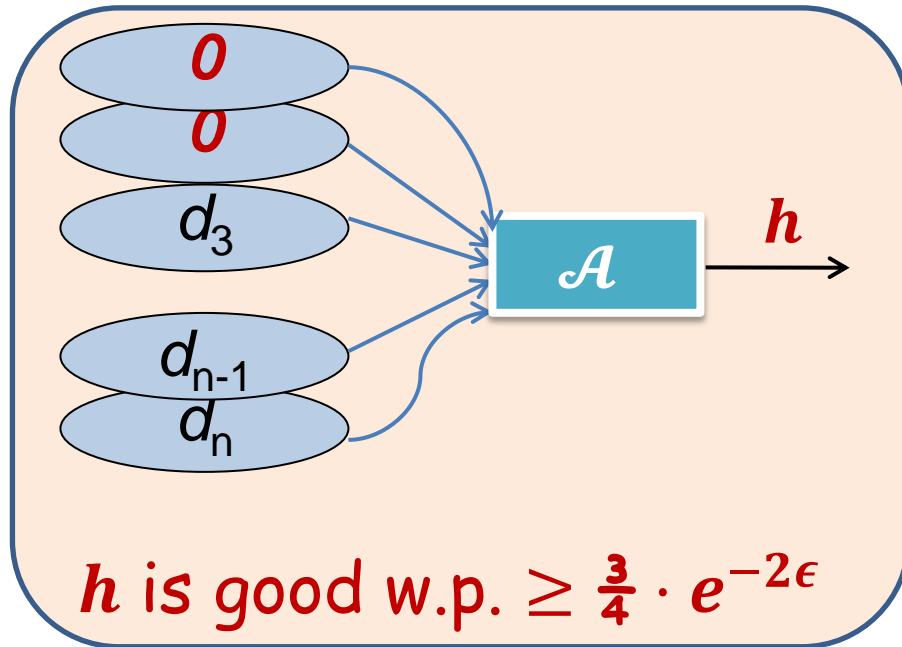
\mathcal{A} : PPAC learner for a concept class \mathcal{C} using n examples.

Correctly sampled data:



Private Learner \implies Prob. Rep.

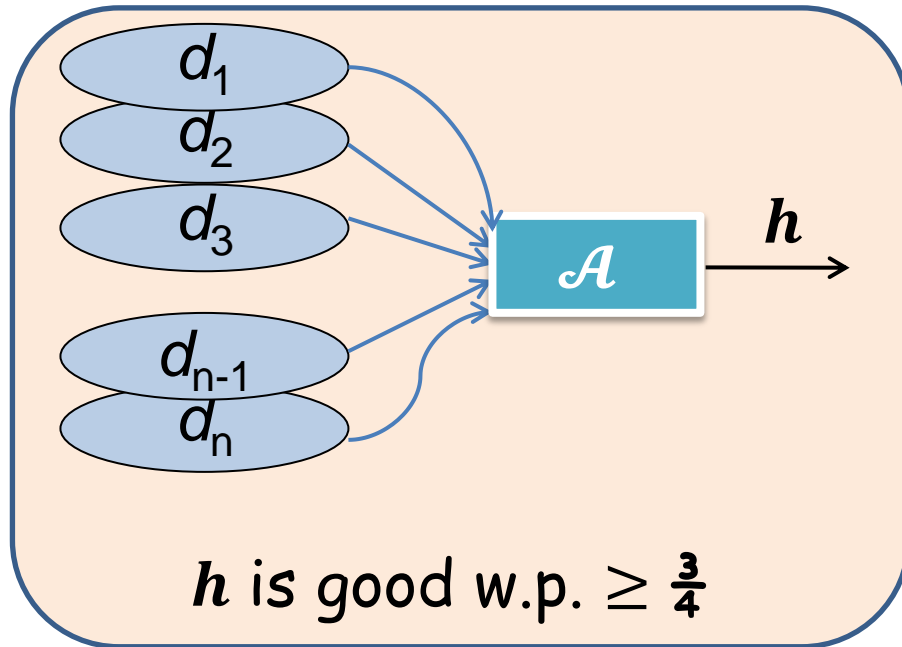
\mathcal{A} : PPAC learner for a concept class \mathcal{C} using n examples.



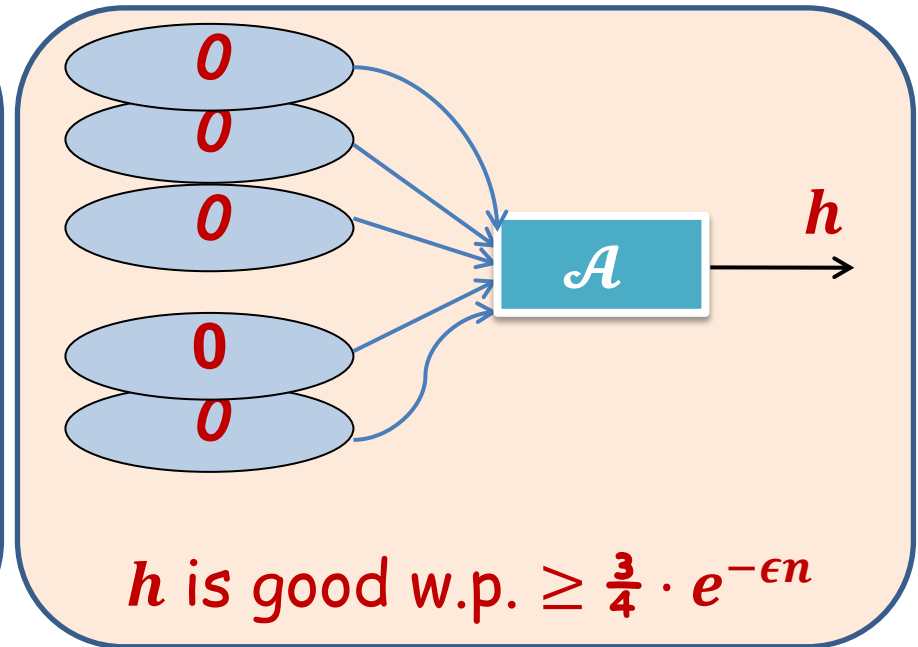
Private Learner \implies Prob. Rep.

\mathcal{A} : PPAC learner for a concept class \mathcal{C} using n examples.

Correctly sampled data:



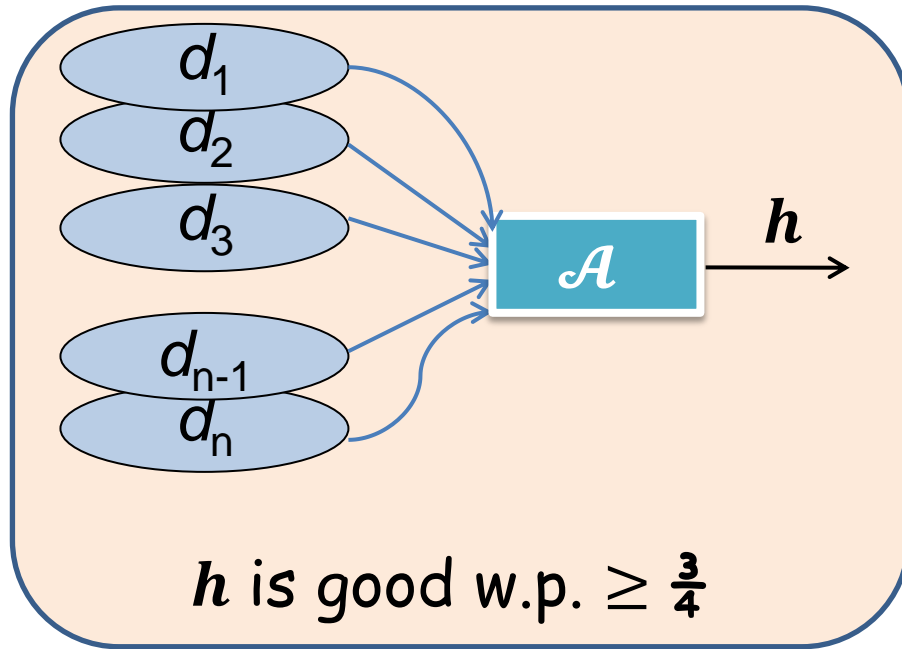
Arbitrary data:



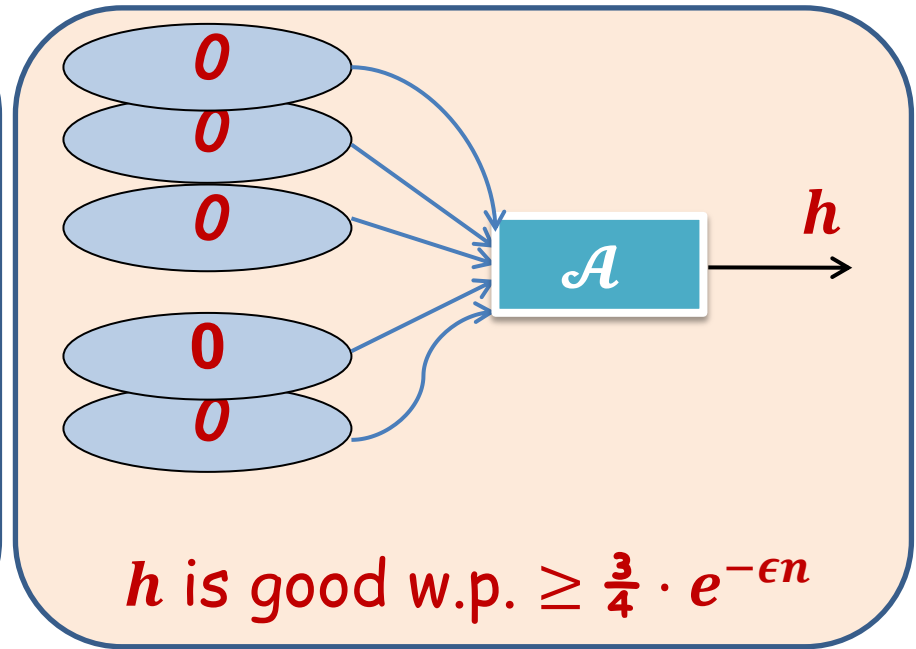
Private Learner \implies Prob. Rep.

\mathcal{A} : PPAC learner for a concept class \mathcal{C} using n examples.

Correctly sampled data:



Arbitrary data:



If we run $\mathcal{A}(0)$ for $O(e^{\epsilon n})$ times, we get a list of hypotheses s.t. with probability $> \frac{3}{4}$ contains a good hypothesis.

- Repeat to get enough lists to create a probabilistic representation

A Tight Characterization

$\Theta(\text{RepDim}(\mathcal{C}))$ samples are necessary and sufficient for privately learning \mathcal{C} .

- Analogous to the VC dimension for non-private learners.

Prob. Reb vs. Det. Rep.

Deterministic Representation: A hypothesis class \mathcal{H} s.t. for every $c \in \mathcal{C}$ and \mathcal{D} , there exists a hypothesis $h_0 \in \mathcal{H}$ s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.

Probabilistic Representation:

A list of hypothesis classes $\{\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_r\}$ s.t. for every $c \in \mathcal{C}$ and \mathcal{D} :
w.p. $\frac{3}{4}$, a randomly chosen \mathcal{H}_i contains an h_0 s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.

Prob. Reb vs. Det. Rep.

Deterministic Representation: A hypothesis class \mathcal{H} s.t. for every $c \in \mathcal{C}$ and \mathcal{D} , there exists a hypothesis $h_0 \in \mathcal{H}$ s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.

Probabilistic Representation:

A list of hypothesis classes $\{\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_r\}$ s.t. for every $c \in \mathcal{C}$ and \mathcal{D} : w.p. $\frac{3}{4}$, a randomly chosen \mathcal{H}_i contains an h_0 s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.

Naïve construction: given $\text{Rep} = \{\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_r\}$ the union of $\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_r$ is a deterministic representation.

- Size grows by $\ln r$.

Prob. Reb vs. Det. Rep.

Deterministic Representation: A hypothesis class \mathcal{H} s.t. for every $c \in \mathcal{C}$ and \mathcal{D} , there exists a hypothesis $h_0 \in \mathcal{H}$ s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.

Probabilistic Representation:

A list of hypothesis classes $\{\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_r\}$ s.t. for every $c \in \mathcal{C}$ and \mathcal{D} : w.p. $\frac{3}{4}$, a randomly chosen \mathcal{H}_i contains an h_0 s.t. $\text{error}_{\mathcal{D}}(h_0) \leq \frac{1}{4}$.

Naïve construction: given $\text{Rep} = \{\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_r\}$ the union of $\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_r$ is a deterministic representation.

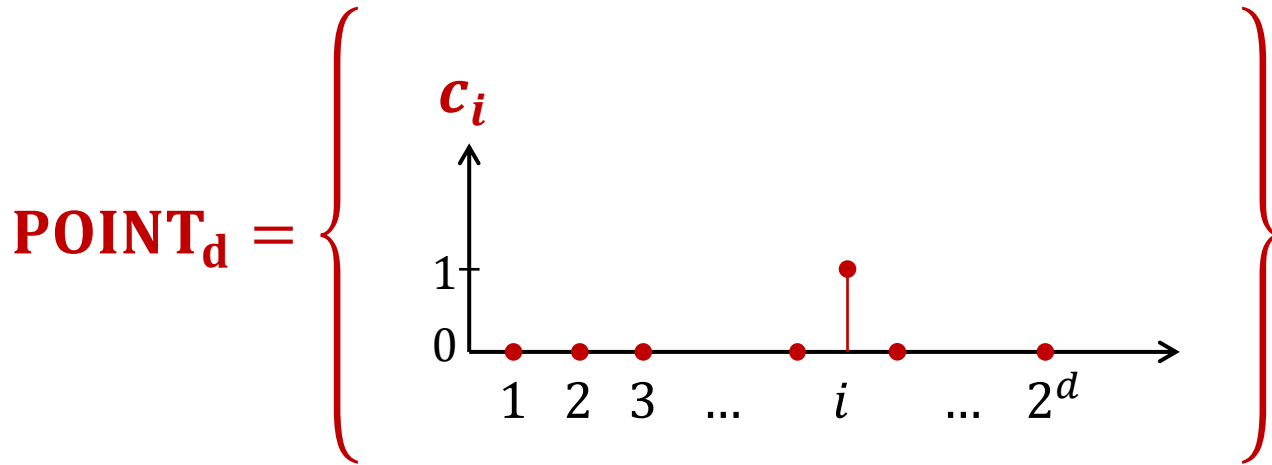
- Size grows by $\ln r$.

Theorem: For every concept class C over the domain $\{0, 1\}^d$ there exists a deterministic representation of C of size $O(\text{RepDim}(C) + \ln d)$

- Non-constructive proof via the probabilistic method.

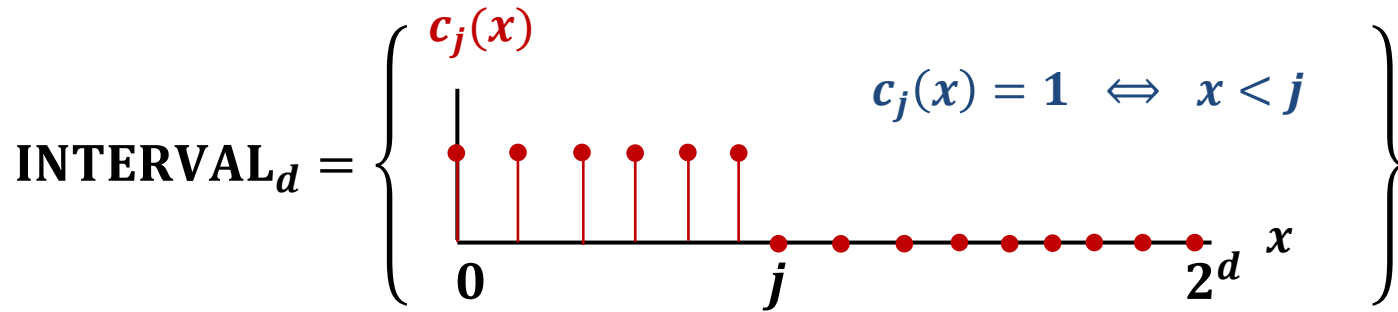
Question

Clearly, $\text{RepDim}(C) \geq \text{VC}(C)$. Is there a concept class C for which $\text{RepDim}(C) \gg \text{VC}(C)$?

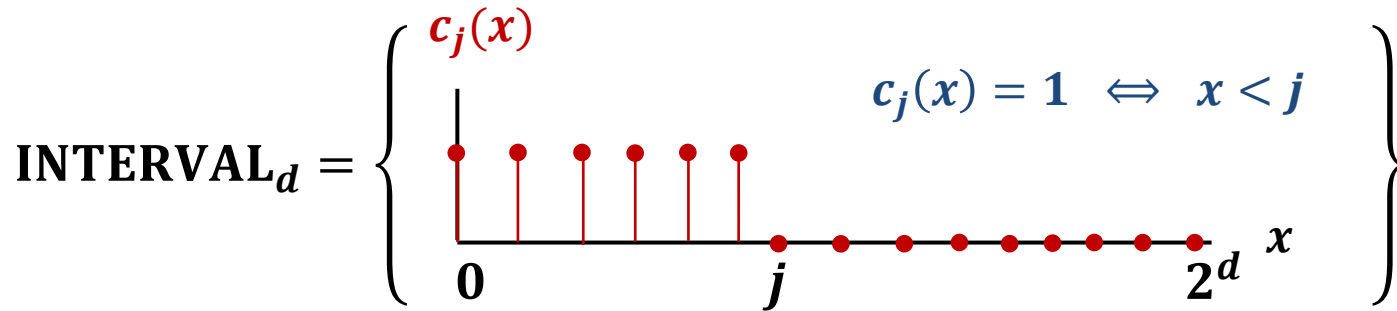


$$\text{RepDim}(\text{POINT}_d) = \text{VC}(\text{POINT}_d) = O(1).$$

Intervals



Intervals



Facts about the sample complexity:

- Non-privately: $O(1)$
- [BKN 2010] ϵ -private proper-learner: $\Omega(d)$
- ϵ -private improper-learner: ???

Approx. D.P. to the Rescue

Dwork, McSherry, Nissim, Smith 2006

Dwork, Kenthapadi, McSherry, Mironov, Naor 2006

A (rand) algorithm \mathcal{A} is (ϵ, δ) differentially private if for all neighboring databases S_1, S_2 and for all sets of outputs F :

$$\Pr[\mathcal{A}(S_1) \in F] \leq e^\epsilon \cdot \Pr[\mathcal{A}(S_2) \in F] + \delta$$

Approx. D.P. to the Rescue

Dwork, McSherry, Nissim, Smith 2006

Dwork, Kenthapadi, McSherry, Mironov, Naor 2006

A (rand) algorithm \mathcal{A} is (ϵ, δ) differentially private if for all neighboring databases S_1, S_2 and for all sets of outputs F :

$$\Pr[\mathcal{A}(S_1) \in F] \leq e^\epsilon \cdot \Pr[\mathcal{A}(S_2) \in F] + \delta$$

We Show: (RANDOM'13)

We show (ϵ, δ) -private proper-learner with sample complexity $2^{O(\log^* d)}$.

Summary and Open Problems

What we saw:

- A rich picture for private learners
 - Even for simple classes such as Points and Intervals.
- Sample complexity of private learners is characterized by the representation dimension.
 - Similar characterization (not as clean) for other tasks.

Summary and Open Problems

What we saw:

- A rich picture for private learners
 - Even for simple classes such as Points and Intervals.
- Sample complexity of private learners is characterized by the representation dimension.
 - Similar characterization (not as clean) for other tasks.

Open problems:

- Tools for lower bounding the representation dimension
- Clearly, $\text{RepDim}(C) \geq \text{VC}(C)$.
Is there a concept class C for which $\text{Repdim}(C) \gg \text{VC}(C)$?
- Characterization of the sample complexity for (ϵ, δ) -differential privacy.